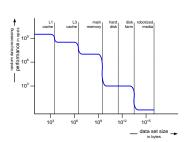
Modeling and Analysis Fundamentals of Technology

-



Gerrit Muller

University of South-Eastern Norway-NISE Hasbergsvei 36 P.O. Box 235, NO-3603 Kongsberg Norway

gaudisite@gmail.com

Abstract

This presentation shows fundamental elements for models that are ICT-technology related. Basic hardware functions are discussed: storage, communication and computing with fundamental characteristics, such as throughput, latency, and capacity. A system is build by layers of software on top of hardware. The problem statement is how to reason about system properties, when the system consists of many layers of hardware and software.

Distribution

This article or presentation is written as part of the Gaudí project. The Gaudí project philosophy is to improve by obtaining frequent feedback. Frequent feedback is pursued by an open creation process. This document is published as intermediate or nearly mature version to get feedback. Further distribution is allowed as long as the document remains complete and unchanged.

All Gaudí documents are available at: http://www.gaudisite.nl/

version: 0.5 status: preliminary draft June 21, 2020

1 Introduction

Figure 1 provides an overview of the content. In this article we discuss generic know how of computing technology. We will start with a commonly used decomposition and layering. We provide *figures of merit* for several generic computing functions, such as storage and communication. Finally we discuss caching as example of a technology that is related to storage figures of merit. We will apply the caching in a web shop example, and discuss design considerations.



Figure 1: Overview Content Fundamentals of Technology

When we model technology oriented design questions we often need feasibility answers that are assessed at the level of non functional system requirements. Figure 2 shows a set of potential technology questions and the required answers at system level.

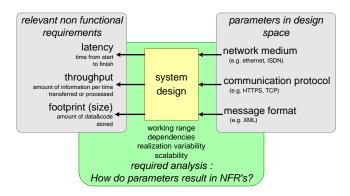


Figure 2: What do We Need to Analyze?

From design point of view we need, for example, information about the working range, dependencies, variability of the actual realization, or scalability.

2 Computing Technology Figures of Merit

In information and communication systems we can distinguish the following generic technology functions:

storage ranging from short term volatile storage to long term persistent storage. Storage technologies range from solid state static memories to optical disks or tapes.

communication between components, subsystems and systems. Technologies range from local interconnects and busses to distributed networks.

processing of data, ranging from simple control, to presentation to compute intensive operations such as 3D rendering or data mining. Technologies range from general purpose CPUs to dedicated I/O or graphics processors.

presentation to human beings, the final interaction point with the human users. Technologies range from small mobile display devices to large "cockpit" like control rooms with many flat panel displays.

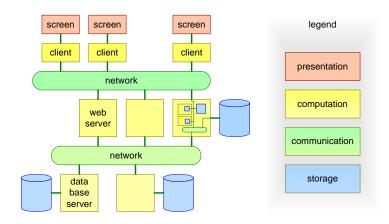


Figure 3: Typical Block Diagram and Typical Resources

Figure 3 shows these four generic technologies in the typical layering of a *Service Oriented Architecture* (SOA). In such an architecture the repositories, the bottom-tier of this figure, are decoupled from the business logic that is being handled in the middle layer, called *web server*. The client tier is the access and interaction layer, which can be highly distributed and heterogeneous.

The four generic technologies are recursively present: within a web-server, for example, communication, storage and processing are present. If we would zoom in further on the CPU itself, then we would again see the same technologies.

		latericy	capacity
processor cache	L1 cache L2 cache L3 cache	sub ns	n kB n MB
fast volatile	main memory	tens ns	n GB
persistent	disks disk arrays disk farms	ms	n*100 GB n*10 TB
archival	robotized optical media tape	a >s	n PB

Lx.

Figure 4: Hierarchy of Storage Technology Figures of Merit

For every generic technology we can provide *figures of merit* for several characteristics. Figure 4 shows a table with different storage technologies. The table provides typical data for latency and storage capacity. Very fast storage technologies tend to have a small capacity. For example, L1 caches, static memory as part of the CPU chip, run typically at processor speeds of several GHz, but their capacity is limited to several kilobytes. The much higher capacity main memory, solid state dynamic RAM, is much slower, but provides Gigabytes of memory. Non solid state memories use block access: data is transferred in chunks of many kilobytes. The consequence is that the access time for a single byte of information gets much longer, milliseconds for hard disks. When mechanical constructions are needed to transport physical media, such as robot arms for optical media, then the access time gets dominated by the physical transport times.

Figure 5 shows the same storage figures of merit in a 2-dimensional graph. The horizontal axis shows the capacity or the maximum data set size that we can store. The vertical axis shows the latency if we axis a single byte of information in the data set in a random order. Note that both axes are shown as a logarithmic scale, both axes cover a dynamic range of many orders of magnitude! The resulting graph shows a rather non-linear behavior with step-like transitions. We can access data very fast up to several kilobytes; the access time increases significantly when we exceed the L1 cache capacity. This effect repeats itself for every technology transition.

The communication figures of merit are shown in the same way in Figure 6. In this table we show *latency*, *frequency* and *distance* as critical characteristics. The latency and the distance have a similar relationship as latency and capacity for storage: longer distance capabilities result in longer latencies. The frequency

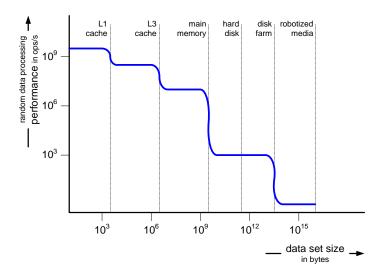


Figure 5: Performance as Function of Data Set Size

behavior, which relates directly to the transfer capacity, is different. On chip very high frequencies can be realized. Off chip and on the printed circuit board these high frequencies are much more difficult and costly. When we go to the long-distance networks optical technologies are being used, with very high frequencies.

3 Caching in Web Shop Example

The speed differences in storage and communication often result in the use of a cache design pattern. The cache is a local fast storage, where frequently used data is stored to prevent repeated slow accesses to slow storage media. Figure 7 shows that this caching pattern is applied at many levels within a system, for example:

network layer cache to avoid network latencies for distributed data. Many communication protocol stacks, such as http, have local caches.

file cache as part of the operating system. The file cache caches the stored data itself as well as directory information in main memory to speed up many file operations.

application cache application programs have dedicated caches based on application know how.

L1, L2, L3 memory caches A multi-level cache to bridge the speed gap between on-chip speed and off chip dynamic memory.

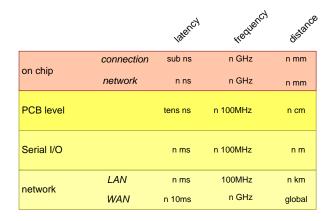


Figure 6: Communication Technology Figures of Merit

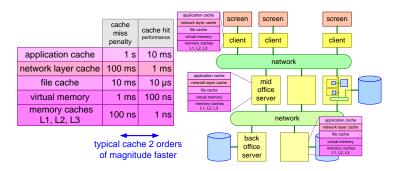


Figure 7: Multiple Layers of Caching

virtual memory where the physical main memory is cache for the much slower virtual memory that resides mostly on the hard disk.

Note that in the 3-tier SLA approach these caches are present in most of the tiers. In Figure 8 we analyze the introduction of caches somewhat more. At the left hand side we show that *long latencies of storage and communication, communication overhead*, and *resource intensive processing* are the main reasons to introduce caching. In the background the project needs for performance and cost are seen as driving factors. Potential performance problems could also be solved by over-dimensioning, however this might conflict with the cost constraints on the project.

The design translates these performance reasons into a number of design choices:

frequently used subset enable the implementation to store this subset in the low capacity, but faster type of memory.

fast storage relates immediately to low latency of the storage itself

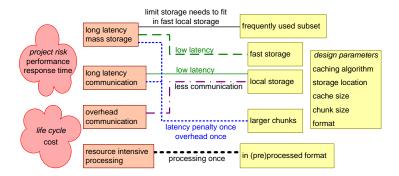


Figure 8: Why Caching?

local storage gives low latency for the communication with the storage (sub)systemlarger chunks reduces the number of times that storage or communication latency occurs and reduces the overhead.

cache in (pre)processed format to reduce processing latency and overhead

These design choices again translate in a number of design parameters:

- caching algorithm
- storage location
- cache size
- · chunk size
- format

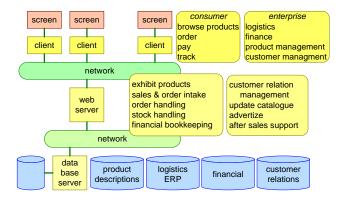


Figure 9: Example Web Shop

As an example of caching we look at a web shop, as shown in Figure 9. Customers at client level should be able to browse the product catalogue, to order products, to pay, and to track the progress of the order. Other stakeholders at client level have logistics functions, financial functions, and can do product and customer management. The web server layer provides the logic for the exhibition of products, the sales and order intake, the order handling, the stock handling, and the financial bookkeeping. Also at the web server layer is the logic for customer relation management, the update of the product catalogue, the advertisements, and the after sales support. The data base layer has repositories for product descriptions, logistics and resource planning, customer relations, and financial information.

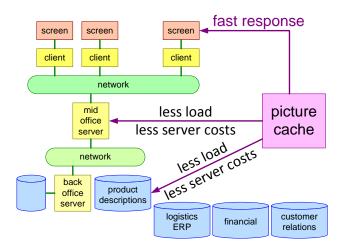


Figure 10: Impact of Picture Cache

We will zoom in on the product browsing by the customers. During this browsing customers can see pictures of the products in the catalogue. The originals of these pictures reside in the product catalogue repository in the data base layer. The web server determines when and how to show products for customers. The actual pictures are shown to many customers, who are distributed widely over the country.

The customers expect a fast response when browsing. Slow response may result in loss of customer attention and hence may cause a reduced sales. A picture cache at the web server level decreases the load at web server level, and at the same time improves the response time for customer browsing. It also reduces the server load of the data base.

So far, the caching appears to be a no-brainer: improved response, reduces server loads, what more do we want? However, Figure 11 shows the potential risks of caching, caused mostly by increased complexity and decreased transparency. These risks are:

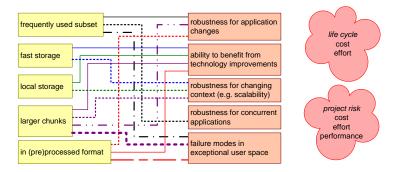


Figure 11: Risks of Caching

- The robustness for application changes may decrease, because the assumptions are not true anymore.
- The design becomes specific for this technology, impacting the ability to benefit from technology improvements.
- The robustness for changing context (e.g. scalability) is reduced
- The design is not robust for concurrent applications
- Failure modes in exceptional user space may occur

All of these technical risks translate in project risks in terms of cost, effort and performance.

4 Summary

Conclusions

Technology characteristics can be discontinuous

Caches are an example to work around discontinuities

Caches introduce complexity and decrease transparancy

Techniques, Models, Heuristics of this module

Generic block diagram: Presentation, Computation, Communication and Storage

Figures of merit

Local reasoning (e.g. cache example)

Figure 12: Summary

Figure 12 shows a summary of this paper. We showed a generic block diagram with *Presentation*, *Computation*, *Communication* and *Storage* as generic computing technologies. Technology characteristics of these generic technologies have discontiuous characteristics. At the transition from one type of technology to another type of technology a steep transition of characteristics takes place. We have provided *figures of merit* for several technologies. Caches are an example to work around these discontinuities. However, caches introduce complexity and decrease the transparancy of the design. We have applied local reasoning graphs to discuss the reasons of introduction of caches and the related design parameters. later we applied the same type of graph to discuss potential risks caused by the increased complexity and decreased transparancy.

References

[1] Gerrit Muller. The system architecture homepage. http://www.gaudisite.nl/index.html, 1999.

History

Version: 0.5, date: 23 February 2007 changed by: Gerrit Muller

updated summary

- Version: 0.4, date: 22 February 2007 changed by: Gerrit Muller

 added text

 Version: 0.3, date: 2 January 2007 changed by: Gerrit Muller

 added sheet about cache levels

 added whe shop example

 removed layering diagram

 Version: 0.2, date: 17 November 2006 changed by: Gerrit Muller

 completely reworked characteristics slide

 many small improvements

 changed status to preliminary draft

 Version: 0.1, date: 10 November 2006 changed by: Gerrit Muller

 added introductory slide

 Version: 0, date: 6 November 2006 changed by: Gerrit Muller

 Created, no changelog yet